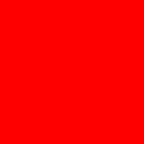




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Project Coin: Language Evolution in the Open

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coin, *n.* A piece of small change
coin, *v.* To create new language

Outline

- Background
- Overview of new language features
- Demo of features in NetBeans
- Developing the features
- Q & A

From *Evolving the Java™ Language*, JavaOne 2005

- **Java Language Principles**
 - Reading is more important than writing
 - Code should be a joy to read
 - The language should not hide what is happening
 - Code should do what it seems to do
 - Simplicity matters
 - A clear semantic model greatly boosts readability
 - Every “good” feature adds more “bad” weight
 - Sometimes it is best to leave things out

Evolving the Java™ Language, JavaOne 2005, cont.

- **One language: with same meaning everywhere**
- **We will evolve the Java Language but cautiously, with a long-term view**
 - we want Java to be around in 2030
 - we can't take a slash-and-burn approach
 - “first do no harm”
- **We will add a few selected features periodically**
 - aimed at developer productivity
 - while preserving clarity and simplicity

Project Coin Today

- OpenJDK Project:
<http://openjdk.java.net/projects/coin/>
- JSR 334
<http://www.jcp.org/en/jsr/detail?id=334>
- To try out Coin features:
 - Build from source yourself!
<http://hg.openjdk.java.net/jdk7/jdk7>
 - Use an IcedTea 7 build
 - Download proprietary JDK 7 binary snapshot builds from
<http://jdk7.dev.java.net/>
- Current JDK 7 GA is July 2011

Project Coin Features in JDK 7

- Binary literals and underscores in literals
- Strings in switch
- Varargs warnings
- Diamond
- Multi-catch and more precise rethrow
- **try-with-resources**
(formerly known as Automatic Resource Management, ARM)

Project Coin Tomorrow?

- Collections support?
 - Collection literals?
 - Support for [] access?
- Large arrays?
- Unsigned integer literals?
- Multi-line strings??
- Your favorite feature???

Coin Constraints

- *Small* language changes
 - Small in specification, implementation, testing
 - No new keywords!
 - Wary of type system changes
 - No JVM changes
- Coordinate with larger language changes
 - Project Lambda
 - Modularity
- One language, one **javac**
 - Interplay and interactions



The Features

A Java Riddle

What is special about the `int` value

1346704470



What is special about the `int` value

1 346 704 470



What is special about the `int` value

`0b01010000010001010001010001010110`

What is special about the `int` value

`0b0101_0000_0100_0101_0001_0100_0101_0110`

What is special about the `int` value

`0b0101_0000_0100_0101_0001_0100_0101_0110`

From the lsb, bit positions set are

2, 3, 5, 7, 11, 13, 17, 19...



What is special about the `int` value

0b0101_0000_0100_0101_0001_0100_0101_0110

From the lsb, bit positions set are

2, 3, 5, 7, 11, 13, 17, 19...

The bits set are the prime bit positions!

Strings in Switch

- When do you use a switch statement?
 - Many alternatives
- Case labels include
 - Integral *constants*
 - Enum constants
- But strings can be constants too!

```
int monthNameToDays(String s, int year) {
    if(s.equals("April") ||
        s.equals("June") ||
        s.equals("September") ||
        s.equals("November"))
        return 30;
    if(s.equals("January") ||
        s.equals("March") ||
        s.equals("May") ||
        s.equals("July") ||
        s.equals("August") ||
        s.equals("December"))
        return 31;
    if(s.equals("February"))
        ...
    else
        ...
}
}
```

```
int monthNameToDays(String s, int year) {
    if(s == "April" ||
        s == "June" ||
        s == "September" ||
        s == "November")
        return 30;
    if(s == "January" ||
        s == "March" ||
        s == "May" ||
        s == "July" ||
        s == "August" ||
        s == "December")
        return 31;
    if(s == "February")
        ...
    else
        ...
}
}
```

```
int monthNameToDays(String s, int year) {
    switch(s) {
        case "April":
        case "June":
        case "September":
        case "November":
            return 30;
        case "January":
        case "March":
        case "May":
        case "July":
        case "August":
        case "December":
            return 31;
        case "February":
            ...
        default
            ...
    }
}
```

```
int monthNameToDays(String s, int year) {
    switch(s) {
        case "April":           case "June":
        case "September":       case "November":
            return 30;
        case "January":         case "March":
        case "May":              case "July":
        case "August":           case "December":
            return 31;
        case "February":
            ...
        default
            ...
    }
}
```

Varargs warnings

- Is anything wrong with calling
 - `Arrays.asList(T... a)`
 - `Collections.addAll(Collection<? super T> c, T... elements)`
 - `EnumSet.of(E first, E... rest)`
- No!

```
class Test {
    public static void main(String... args) {
        List<List<String>> monthsInTwoLanguages =
            Arrays.asList(Arrays.asList("January",
                                        "February"),
                          Arrays.asList("Gennaio",
                                        "Febbraio"));
    }
}
```



```
class Test {
    public static void main(String... args) {
        List<List<String>> monthsInTwoLanguages =
            Arrays.asList(Arrays.asList("January",
                                     "February"),
                        Arrays.asList("Gennaio",
                                     "Febbraio"));
    }
}
```

```
Test.java:7: warning:
[unchecked] unchecked generic array creation
for varargs parameter of type List<String>[]
    Arrays.asList(Arrays.asList("January",
                               ^
1 warning
```

Heap Pollution – JLSv3 4.12.2.1

- For example, a variable of type `List<String>[]` might point to an array of Lists where the Lists did not contain strings
- Reports possible locations of `ClassCastException` at runtime
- A consequence of erasure and lack of reification

```
class Test {
    public static void main(String... args) {
        List<List<String>> monthsInTwoLanguages =
            Arrays.asList(Arrays.asList("January",
                                     "February"),
                        Arrays.asList("Gennaio",
                                     "Febbraio"));
    }
}
```

```
Test.java:7: warning:
[unchecked] unchecked generic array creation
for varargs parameter of type List<String>[]
    Arrays.asList(Arrays.asList("January",
                             ^
1 warning
```

But nothing bad happens!

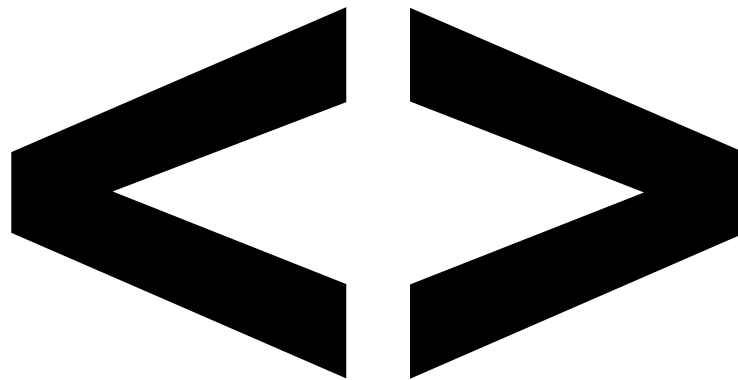
- Arrays created by the compiler for varargs are properly formed
- Well-behaved methods just iterate over the elements
- Unfriendly to warn at every call site
- *Declaration* is problematic

Varargs Warnings Revised

- New mandatory compiler warning at suspect varargs method declarations
- By applying an annotation at the declaration, warnings at the declaration *and call sites* can be suppressed
- New “**@SafeVarargs**” annotation in `java.lang`
 - Compiler will trust, may verify
 - Warnings or errors if annotation applied improperly
 - JDK libraries retrofitted where appropriate, :
“Project Coin: Safe Varargs in JDK Libraries,”
http://blogs.sun.com/darcy/entry/project_coin_safe_vararg_libraries

```
class Test {
    public static void main(String... args) {
        List<List<String>> monthsInTwoLanguages =
            {{ "January", "February" },
             { "Gennaio", "Febbraio" }};
    }
}
```

A possible future with Collection literals.



Pre-generics

```
List list =  
    new ArrayList();
```


With Generics

```
List<String> list =  
    new ArrayList<String> ();
```

With diamond

```
List<String> list =  
    new ArrayList<>();
```

```
List<List<List<List<List<String>>>>> list =  
    new ArrayList<List<List<List<List<String>>>>> ();
```

```
List<List<List<List<List<String>>>>> list =  
    new ArrayList<>();
```

Mining diamonds in the libraries

- Systematic effort to retrofit diamond into JDK libraries
- Potential diamond locations found using javac & NetBeans technologies
- Distribution of diamonds:
 - field or local variable initializer (61%)
 - right-hand side of assignment statement (33%)
 - method argument (4%)
 - return statement (2%)

Multi-catch with More Precise Rethrow

```
try {
    // Reflective operations calling Class.forName,
    // Class.newInstance, Class.getMethod,
    // Method.invoke, etc.
} catch (ClassNotFoundException cnfe) {
    log(cnfe);
    throw cnfe;
} catch (InstantiationException ie) {
    log(ie);
    throw ie;
} catch (NoSuchMethodException nsme) {
    log(nsme);
    throw nsme;
} catch (InvocationTargetException ite) {
    log(ite);
    throw ite;
}
```

A tempting, but troublesome alternative

```
try {  
    // Reflective operations calling Class.forName,  
    // Class.newInstance, Class.getMethod,  
    // Method.invoke, etc.  
} catch (Exception e) {  
    log(e);  
    throw e;  
}
```

Exception by-catch

```
try {  
    // Reflective operations calling Class.forName,  
    // Class.newInstance, Class.getMethod,  
    // Method.invoke, etc.  
} catch (Exception e) {  
    log(e);  
    throw e;  
}
```

**Catches both checked
and unchecked exceptions**

Reduced by-catch

```
try {  
    // Reflective operations calling Class.forName,  
    // Class.newInstance, Class.getMethod,  
    // Method.invoke, etc.  
} catch (RuntimeException e) {  
    ...  
} catch (Exception e) {  
    log(e);  
    throw e;  
}
```

Better.

Multi-catch

```
try {  
    // Reflective operations calling Class.forName,  
    // Class.newInstance, Class.getMethod,  
    // Method.invoke, etc.  
} catch (ClassNotFoundException |  
         InstantiationException |  
         NoSuchMethodException |  
         InvocationTargetException e) {  
    log(e);  
    throw e;  
}
```

More Precise Rethrow

```
try {  
    // Reflective operations calling Class.forName,  
    // Class.newInstance, Class.getMethod,  
    // Method.invoke, etc.  
} catch (ClassNotFoundException |  
        InstantiationException |  
        NoSuchMethodException |  
        InvocationTargetException e) {  
    log(e);  
    throw e;  
}
```

More More Precise Rethrow

```
try {  
    // Reflective operations calling Class.forName,  
    // Class.newInstance, Class.getMethod,  
    // Method.invoke, etc.  
} catch (ReflectiveOperationException e) {  
    log(e);  
    throw e; // Means ClassNotFoundException or ...  
}
```

Still More More Precise Rethrow

```
void foo() throws ClassNotFoundException ...
try {
    // Reflective operations calling Class.forName,
    // Class.newInstance, Class.getMethod,
    // Method.invoke, etc.
} catch (ReflectiveOperationException e) {
    log(e);
    throw e; // Means ClassNotFoundException or ...
}
```

More precise rethrow

- Under `-source 7`, enabled by default for `final` and *effectively final* catch parameters
- From *quantitative analysis*, $99\frac{44}{100}\%$ of catch parameters are `final` or effectively final
- Changing meaning of `throw`
 - Stops compilation of contrived legal programs, *but*
 - Compilation breakage not observed in practice analyzing 9+ million loc in several dozens projects
- *Disjunctive* catch parameters are implicitly final
 - Eases fuller support for disjunctive types in the future

try-with-resources **(Automatic Resource Management)**

- Let's say you want to copy an input stream to an output stream...

```
InputStream in = new FileInputStream(src);  
OutputStream out = new FileOutputStream(dest);
```

```
byte[] buf = new byte[8192];  
int n;
```

```
while ((n = in.read(buf)) >= 0)  
    out.write(buf, 0, n);
```



```
InputStream in = new FileInputStream(src);  
OutputStream out = new FileOutputStream(dest);
```

```
byte[] buf = new byte[8192];  
int n;
```

```
while ((n = in.read(buf)) >= 0)  
    out.write(buf, 0, n);
```

```
InputStream in = new FileInputStream(src);
OutputStream out = new FileOutputStream(dest);
try {
    byte[] buf = new byte[8192];
    int n;

    while ((n = in.read(buf)) >= 0)
        out.write(buf, 0, n);
} finally {
    out.close();
    in.close();
}
```

```
InputStream in = new FileInputStream(src) ;
try {
    OutputStream out = new FileOutputStream(dest) ;
    try {
        byte[] buf = new byte[8192];
        int n;

        while ((n = in.read(buf)) >= 0)
            out.write(buf, 0, n);
    } finally {
        out.close();
    }
} finally {
    in.close();
}
```

```
InputStream in = new FileInputStream(src);
try {
    OutputStream out = new FileOutputStream(dest) ;
    try {
        byte[] buf = new byte[8192];
        int n;

        while ((n = in.read(buf)) >= 0)
            out.write(buf, 0, n);
    } finally {
        out.close();
    }
} finally {
    in.close();
}
```

**What if an exception
occurs here?**

```
InputStream in = new FileInputStream(src);
try {
    OutputStream out = new FileOutputStream(dest) ;
    try {
        byte[] buf = new byte[8192];
        int n;

        while ((n = in.read(buf)) >= 0)
            out.write(buf, 0, n);
    } finally {
        out.close();
    }
} finally {
    in.close();
}
```

**Can get another
exception here!**

```
InputStream in = new FileInputStream(src);
try {
    OutputStream out = new FileOutputStream(dest) ;
    try {
        byte[] buf = new byte[8192];
        int n;

        while ((n = in.read(buf)) >= 0)
            out.write(buf, 0, n);
    } finally {
        out.close();
    }
} finally {
    in.close();
}
```

**Could even get
a third exception here!**

Considerations

- First exception thrown is most likely to be informative
- Exception from a `close` method should propagate, unless there is already an incoming exception
- Don't want to lose all record of a *suppressed* exception
- The additional code to implement this doesn't fit on a slide anymore

```
InputStream in = new FileInputStream(src);
OutputStream out = new FileOutputStream(dest);

byte[] buf = new byte[8192];
int n;

while ((n = in.read(buf)) >= 0)
    out.write(buf, 0, n);
```



```
try(InputStream in = new FileInputStream(src);
    OutputStream out = new FileOutputStream(dest)) {

    byte[] buf = new byte[8192];
    int n;

    while ((n = in.read(buf)) >= 0)
        out.write(buf, 0, n);
}
```

How sweet it is

- Compiler desugars `try-with-resources` into nested `try-finally` blocks with variables to track exception state
- Suppressed exceptions are recorded for posterity using a new facility of **Throwable**
- API support in JDK 7
 - New superinterface `java.lang.AutoCloseable`
 - All **AutoCloseable** and by extension `java.io.Closeable` types usable with `try-with-resources`
 - JDBC 4.1 retrofitted as **AutoCloseable** too

More informative backtraces

```
java.io.IOException
  at Suppress.write(Suppress.java:19)
  at Suppress.main(Suppress.java:8)
  Suppressed: java.io.IOException
    at Suppress.close(Suppress.java:24)
    at Suppress.main(Suppress.java:9)
  Suppressed: java.io.IOException
    at Suppress.close(Suppress.java:24)
    at Suppress.main(Suppress.java:9)
```

To update your code to use with `try-with-resources`

- All `Closeables` are already useable!
- If a type has a no-arg `public void close()` method, implement `AutoCloseable` or `Closeable` as appropriate
- Use an annotation processor to find types to retrofit:
“Project Coin: Bringing it to a `Close(able)`,”
http://blogs.sun.com/darcy/entry/project_coin_bring_close

Demo

NetBeans daily build:

<http://bits.netbeans.org/download/trunk/nightly/latest/>

IDEA (work in progress):

<http://confluence.jetbrains.net/display/IDEADEV/IDEA+X+EAP>

**Eclipse: No bits with Project Coin support to download,
but EclipseCon is next month!**

Developing the features

- Straightforward to use, less obvious to develop!

So you want to change the language...

- Update the Java Language Spec.
- Compiler Implementation
- Essential library support
- Write tests
- Update the JVM Spec.
- Future language evolution
- Update the JVM and class file tools
- Update JNI
- Update the reflective APIs
- Update serialization
- Update javadoc output
- Kinds of compatibility

Updating the Java Language Specification

- Syntax
- Type system
- Method resolution
- Flow analysis, e.g. definite assignment
- Memory model
- Total length of JLSv3: 647 pages
 - Chapter on Lexical structure ends on page 32
 - Syntax chapter is 12 pages
 - Syntax is less than 6% of the JLS!

Writing language change unit/regression tests

- “Writing javac regression and unit tests for new language features,”
http://blogs.sun.com/darcy/entry/javac_regression_tests
- Negative tests:
 - Invalid source files are rejected with expected error messages referencing the proper source locations
- Positive tests:
 - Valid source is compiled.
 - Proper modeling of the new language construct.
 - Resulting class files are structurally well-formed.
 - Resulting class files follow compiler-specific idioms.
 - Resulting class files run have correct operational semantics.

Strings in switch specification change

JLS §14.11 The switch Statement

“The type of [the switch] *Expression* must be **char**, **byte**, **short**, **int**, **Character**, **Byte**, **Short**, **Integer**, **String**, or an enum type (§8.9), or a compile-time error occurs.”

Strings in switch Project Coin proposal form

PROJECT COIN NAME: LANGUAGE CHANGE PROPOSAL NUMBER:
AUTHOR(S): Joseph D. Darcy

OVERVIEW

Provide two sentence or shorter description of these five aspects of the feature:

FEATURE SUMMARY: Should describe it as a summary in a language tutorial.

ADD THE MERITS: switch covering values, analogies to the existing ability to switch on values of the primitive types.

MAIN ADVANTAGE: What makes the proposal a favorable change?

More regular coding patterns can be used for operations selected on the basis of a set of constant string values, thus making of the new construct should be obvious to a developer.

MAIN BENEFIT: Why is the problem better if the proposal is adopted?

Potentially better performance for string constants/paths code.

MAIN DISADVANTAGE: There is always a cost.

Some increased implementation and testing complexity for the compiler.

ALTERNATIVES: Consider benefits and disadvantages of each way without a language change?

No, chosen if the other tests for string equality are generally expensive and introducing an means for it work like constants, one per string value of interest, would add another type to a program without justification.

EXAMPLES

Show the code snippets (SAMPLES). Show the simplest possible program utilizing the new feature.

```
String s = ...
switch (s)
{
    case "foo":
        processFoo(s);
        break;
}
```

ADVANCED EXAMPLES: Show advanced usage of the feature.

```
String s = ...
switch (s)
{
    case "foo":
        processFoo(s);
        // fall through
    case "bar":
        processBar(s);
        break;
}
```

DISCUSS

DISCUSSION: Describe how the proposal affects things like memory, type system, and meaning of expressions and statements in the new Programming language as well as any other known impacts.

The level of grammar is unchanged. String is added to the set of types subject for a switch statement in 5.3 of section 14.2.1. Since Strings are already included in the definition of constant expressions, in 5.3 of section 5.29, that switch label production does not need to be augmented.

The existing switch rules in 14.2.1 are not duplicated, at most one default, no null, and, etc. is applied to Strings as well. The type system is unchanged. The definition of constant expressions is not changed, as of section 6.2.9, is unchanged as well.

IMPLEMENTATION: How would the feature be compiled to class files?

The way to support this change would have to augment the VM's local/global instruction to operate on String values, however, that approach is not recommended or necessary. It would be possible to translate the switch to equivalent if-else code, but that would require some non-trivially complex code which is reportedly expensive in code size, but would occur on a predictable first

Integer (or long) function value, computed from the string. The next natural choice for this function is String.hashCode(), but other functions could also be used either alone or in conjunction with hashCode(). The specification of String.hashCode() is a constant behavior (this point). If all handling labels have an identical length, String.length() could be used instead of hashCode(). Given a String, equal() check will be needed to verify the constant string's identity in addition to the evaluation of the existing function because several printing inputs could evaluate to the same output.

A single case label, a single case label with a default, and two case labels can be specified and not equal checks, without function evaluations. If there are collisions in String.hashCode() over the set of constant labels in a switch block, different functions without collisions, on the set of inputs should be used (for example (long)hashCode() & ^ s.length()) in another candidate function.

Here are suggestions to currently legal (as a source for the two examples above where the default hashCode case code collides):

```
// simple example
if (s.equals("foo")) { // case 99% if it coll
    processFoo(s);
}
```

```
// Advanced example
{ // new scope for synthetic variables
    local var $label_default = fall on
    local var $fallthrough = fall on
    $default_label {
        switch (hashCode()) { // case 99% if it coll
            case 42424242: { // case "foo" hashCode()
                $label_default = true;
                break $default_label;
            }
            processFoo(s);
            $fallthrough = true;
            case 123456789: { // "bar" hashCode()
                if (fallthrough && s.equals("foo")) {
                    $label_default = true;
                    break $default_label;
                }
                processFoo(s);
            }
            $fallthrough = true;
            case 987654321: { // "bar" hashCode()
                if (fallthrough && s.equals("bar")) {
                    $label_default = true;
                    break $default_label;
                }
                processFoo(s);
            }
            case 918765: { // "bar" hashCode()
                if (s.equals("bar")) {
                    $label_default = true;
                    break $default_label;
                }
                processFoo(s);
                $fallthrough = true;
            }
            default:
                $label_default = true;
                break $default_label;
            }
        }
        if ($label_default)
            processDefault(s);
    }
}
```

In the advanced example, the local "fall through" variable is needed to track whether a fall through has occurred on the string equality checks can be bypassed. If there are no fall throughs, this variable can be removed. Likewise, if there is no default label in the original code, the \$label_default variable is not needed and a simple "break" can be used instead.

In a translation directly to byte code, there are other case variables can be replaced with goto(s), exposed (this is possible as a source of the code):

```
// Advanced example in pseudo (as a working code)
switch (hashCode()) { // case 99% if it coll
    case 42424242: { // case "foo" hashCode()
        if (s.equals("foo"))
            goto $default_label;
        goto $fallthrough_label;
    }
}
```

```
case 123456789: { // "bar" hashCode()
    if (s.equals("bar"))
        goto $fallthrough_label;
        goto $fallthrough_label;
}
case 987654321: { // "bar" hashCode()
    if (s.equals("bar"))
        goto $fallthrough_label;
}
switch (hashCode) { default, and two case labels
    processCodeBlock(s);
    break;
}
case 918765: { // "bar" hashCode()
    if (s.equals("bar"))
        goto $fallthrough_label;
        processCode(s);
}
default:
    $label_default = true;
    processDefault(s);
    break;
}
```

Not at the compilation, a compiler executing via generics, a constant falling through switch, such as java's -Xlint:SwitchOption and -Xlint:SwitchFallThrough, should check carefully on switch statement based on strings.

TOOLS: How can the feature be tested?

Generating a simple and complex set of binary structural and semantic paper machine parts, to our, conditions to test include switch statements with and without fall through, with and without collisions in the hash codes, with and without default labels.

LIBRARY SUPPORT: Are any supporting libraries needed for the feature?

No.

REFLECTIVITY: Do any of the various classically reflective APIs need to be updated? Think of reflective APIs included but not limited to code reflection (as a long class or java.lang.reflect.*), java.lang.model.*, the class API, and PDA.

Do reflective APIs that control statements in the source language need to be updated. None of one reflection, java.lang.model.*, the class API, and 2.5 in class statements; therefore, they are unaffected. The new APIs in java.lang.reflect.*, java.lang.model.*, and PDA, does not need statements, but the new string API for switch statements is general enough to model the old language without any API changes.

OTHER CHANGES: Do any other parts of the platform need to be updated? Possible include but are not limited to type declaration, and output of the javac tool.

No.

NOTATION: Sketch how a code base could be converted, manually or automatically, to use the new feature.

Look for occurrences of "constant string" equals() for (if (s.equals("constant string")) and replace a constant.

COMPATIBILITY:

BACKWARD COMPATIBLE: Are any previously valid programs now invalid? If so, list one.

All existing programs remain valid.

FORWARD COMPATIBLE: How do source code files of earlier platform versions interact with the feature? Can any new code be compiled? Can any new code be compiled?

The semantics of writing class files and java source files, and are unchanged by this feature.

FOR TESTING: Please include a list of any existing test files related to this proposal.

See 2.2 of Using String and Object in Switch case statements. http://bug.sun.com/bug/view_bug.jsp?bug_id=5012862

See: PDR 9802317 (optional).

No prototype at this time.

Strings in switch Project Coin proposal form

PROJECT COIN NAME: LANGUAGE CHANGE PROPOSAL NUMBER: 0
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OVERVIEW:
Provide two sentence or shorter description of these five aspects of the feature:
RETURN'S SUMMARY: Should describe how a summary in a language feature affects the ability to switch control values, and give rise to the resulting ability to switch on values of the given types.
MAIN ADVANTAGE: What makes the proposal a favorable change?
More regular coding patterns can be used for operations effected on the basis of a set of constant string values, thus making of the new construct should be obvious to new developers.
MAIN BENEFIT: Why is the platform better if the proposal is adopted?
Potentially better performance for string-based paths code.
MAIN DISADVANTAGE: There is always a cost.
Some increased implementation and testing complexity for the compiler.
ALTERNATIVES: Can the benefits and advantages be had some way without a language change?

No, should if the new tests for string equality are generally expensive and introducing an overhead for it to switch table constants, one per string value of interest, would add another type to a program without justification.

EXAMPLES:
Show what the code SHOULD BEHAVE. It should be simple, good, and program solving the new feature.

```
String s = ...  
switch (s) {  
case "foo":  
    process(s);  
break;  
}
```

ADVANCED EXAMPLES: Show advanced usage of the feature.

```
String s = ...  
switch (s) {  
case "foo":  
    process(s);  
break;  
case "bar":  
    process(s);  
break;  
default:  
    process(s);  
break;  
}
```

DISCUSSION: Describe how the proposal affects things currently supported, and how it affects new constructs and statements in the Java Programming Language.

The level of generality is discussed in section 3.2.1. Since strings are already included in the definition of constant expressions, in section 3.2.2, their switch table production does not need to be augmented. The existing restrictions in 3.2.2 are no longer relevant, at most one default, default must come at the end of the switch, and the target values is unchanged. The default may now also be a constant table statement, and not restricted to a single generic word.

COMPLIANCE: How would the feature be implemented in the compiler?
The implementation of the feature is a simple matter of adding a new table statement to the definition of constant expressions, in section 3.2.2. The implementation does not need to be augmented. The existing restrictions in 3.2.2 are no longer relevant, at most one default, default must come at the end of the switch, and the target values is unchanged. The default may now also be a constant table statement, and not restricted to a single generic word.

COMPATIBILITY: How would the feature be implemented in the compiler?
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strings (or long function values) required from the string. The next natural choice for this function is String.hashCode(), whose function could also be used either alone or in conjunction with hashCode(). The specific choice of String.hashCode() is a comment for later in the post. A switching table can be defined by using hashCode() and forward index of hashCode(). Instead a String equal check will be needed to verify the constant string is used. It is added to the evaluation of these new function because of printing inputs could not do so to the same result.

A single constant, a single case label with a default, a set of case labels can be specified to not equal checks, without function evaluations. If there are multiple String hashCode calls on the set of case labels in a switch block, different function without collisions (on basis of of equal checks) can be used for example (long.hashCode() && 0xffffL) | (hashCode() <> 0) another candidate function.

Here are examples to correctly legal (as a source for that two examples show where the default has to be used, and where it is not).

```
// simple example  
if ("foo" != "bar") { // case with "if" is not  
    process(s);  
}
```

```
// advanced example  
// new scope for synthetic variables  
boolean tableDefault = false;  
boolean tableThrough = false;  
switch (table) {  
    switch (table) { // case with "if" is not  
        case "foo": // case with "if" is not  
            tableDefault = true;  
            break; tableDefault;  
        }  
        process(s);  
        tableThrough = true;  
        case "foo": // case with "if" is not  
            if (tableThrough && is equal of "foo") {  
                tableDefault = true;  
                break; tableDefault;  
            }  
            process(s);  
            break;  
        case "bar": // case with "if" is not  
            if (is equal of "bar") {  
                tableDefault = true;  
                break; tableDefault;  
            }  
            process(s);  
            break;  
        default:  
            tableDefault = true;  
            break; tableDefault;  
    }  
}
```

In the advanced example, the boolean "tableThrough" variable is used to track whether a tableThrough occurred on the string equality checks can be ignored. If there is no tableThrough, this variable can be ignored. However, if there is no default label in the original code, the tableDefault variable is not needed and a simple "break" can be used instead.

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```
case "foo": // case with "if" is not  
    process(s);  
break;  
case "bar": // case with "if" is not  
    process(s);  
break;  
default:  
    process(s);  
break;
```

Compilation, cont.

For details on compilation, a compiler emitting the generated code falling through to the next case, such as a "fall through" option, and if the default is not "fall through", should check the code, on switch statements based on strings.

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Migration

Compatibility

The semantics of writing the file and log files are not affected by this feature.

References, prototype?

How to make a diamond

- *Type inference* has the compiler figure out types rather than the programmer writing them out
- The type argument for diamond, “<>”, is inferred by the compiler
- Diamond reapplies existing type inference features to infer types parameters in constructor calls
- Similar to inference for generic methods:
`public static <T> List<T> asList(T... a)`

What's in the box?

```
... = new Box<>(42) ;
```


What's in the box?

```
public class Box<T> {  
    private T value;  
  
    public Box(T value) {  
        this.value = value;  
    }  
  
    T getValue() {  
        return value;  
    }  
}
```

Box<> (42) ;

Types on the left...

```
Box<Integer> box
```

```
Box<Number> box
```

```
Box<Object> box
```

```
Box<?> box
```

```
Box<? extends Comparable<?>> box
```

```
...
```

```
... = new Box<>(42) ;
```

Pick a type for the right, but not just any type

```
Box<Integer> box
```

```
Box<Number> box
```

```
Box<Object> box
```

```
Box<?> box
```

```
Box<? extends Comparable<?>> box
```

```
...
```

```
... = new Box<>(42) ;
```

```
Integer
```

```
Number
```

```
Object
```

```
Comparable<?>
```

```
Object & Comparable<? extends ...>
```

```
...
```

Two inference schemes, “Simple” and “Complex”

- Simple, types parameters from:
 - Assignment context (where available)
- Complex, type parameters from:
 - Assignment context (where available) *plus*
 - Actual arguments to the constructor

Simple algorithm

```
Box<? extends Number> b = new Box<>(42)
```

```
Integer  
Number  
Object  
Comparable<?>  
Object & Comparable<? extends...>  
...
```

Complex algorithm

```
Box<? extends Number> b = new Box<>(42)
```

```
Integer  
Number  
Object  
Comparable<?>  
Object & Comparable<? extends...>  
...
```

A distinction with a difference

Simple: `Box<Number> b = new Box<>(42)`

Complex: `Box<Number> b = new Box<>(42)`

Integer

Number

Object

Comparable<?>

Object & Comparable<? extends...>

...

A distinction with a difference

Simple: `Box<Number> b = new Box<>(42)`

Complex: `Box<Number> b = new Box<>(42)`

```
incompatible types
Box<Number> b = new Box<>(42);
                ^
    required: Box<Number>
    found:    Box<Integer>
1 error
```


Method contexts and algorithms

```
void m(Box<Integer> box) {...}
```

Simple:

```
m(new Box<>(42))
```

Complex:

```
m(new Box<>(42))
```

Integer

Number

Object

Comparable<?>

Object & Comparable<? extends...>

...

Method contexts and algorithms

```
void m(Box<Integer> box) {...}
```

Simple:

```
m(new Box<>(42))
```

```
method m cannot be applied to given types;  
{ m(new Box<>(42)); }  
  ^  
  required: Box<Integer>  
  found: Box<Object>  
1 error
```

Com

2))

Language design for the real world

- Sometimes the simple algorithm is more useful,
but
other times the complex algorithm is more useful
- What to do?
 - Is either one any good?
 - How to choose between them?
- Generate some data!
- *Quantitative* language design

Experimental Methodology, Summer 2009

- Find relevant large code bases (millions of lines of code)
 - OpenJDK
 - Tomcat
 - NetBeans
- Create and run *diamond finder*
- *Measure* effectiveness of algorithms
- Interpret results and decide

Per code base

	OpenJDK	Tomcat	NetBeans
Total new's	104,138	6,048	94,768
Generics new's	5,076	153	12,010
Simple Success	4,409	148	10,670
Complex Success	4,533	148	11,085

Analysis

	OpenJDK	Tomcat	NetBeans
Total new's	104,138	6,048	94,768
Generics new's	5,076	153	12,010
Simple Success	4,409	148	10,670
Complex Success	4,533	148	11,085

- Nontrivial fraction of constructor calls are to generic classes
- Of constructor calls to generic classes, in **90%** of cases the type parameters are successfully inferred
 - Simple infers in one 90% subset
 - Complex infers in a slightly different 90% subset
- Therefore, either algorithm would be effective
- Given equal effectiveness, what other criteria to use?

A way ahead to break the tie

- Neither algorithm is always better than the other
- Neither algorithm is a *subset* of the other
 - Picking one algorithm in JDK N and the other in JDK $(N+1)$ would mean that some code that compiled in JDK N would stop compiling in JDK $(N+1)$
- Decision today constrains decisions tomorrow
- Originally integrated the simple algorithm...

A rising tide lifts all boats

- ... later switched to the complex algorithm because
 - The complex algorithm reuses more inference machinery in the spec and implementation
 - More maintainable, implicit bug fixes for free
 - Better evolution properties
- Since the experiment, anticipate beneficial interactions with future inference improvements
 - Target typing in Project Lambda

A surprise: why is this disallowed?

- Using a more sophisticated inference scheme can be problematic sometimes

```
List<?> arg = . . . ;  
new Box<> (arg) ;
```

A surprise: why is this disallowed?

```
List<?> arg = ... ;  
new Box<>(arg) ;
```

```
cannot infer type arguments for Box<>;  
new Box<>(arg) ;
```

^

```
reason: type argument List<CAP#1>  
inferred for Box<> is not allowed in this context
```

How is \mathbb{T} inferred?

```
<T> List<T> asList(T... t) {...}
```

```
List<?> arg = ...;
```

```
Arrays.asList(arg);
```

Don't mistreat captured types!

```
<T> List<T> asList(T... t)
```

```
List<?> arg = ...;
```

```
Arrays.asList(arg);
```

```
T == List<#capture of ?>
```

How to break a diamond

```
List<?> arg = ... ;  
new Box<>(arg) ;
```

Not just a copy...

```
List<?> arg = ... ;  
new Box<List<?>> (arg) ;
```

Not just a copy...

```
List<?> arg = ... ;  
new Box<List<capture of ?>> (arg) ;
```

Even worse...

```
List<?> arg = ...;  
new Box<List<capture of ?>>(arg) { ... };
```


How does this get compiled?

```
List<?> arg = ... ;  
new a$1 (arg) ;
```

Anonymous classes translate into a new class file with a full set of attributes.

```
class a$1 extends Box<List<capture of ?>> { ... }
```

How does this get compiled?

This signature cannot be represented in the class file!
Problematic for core reflection and separate compilation

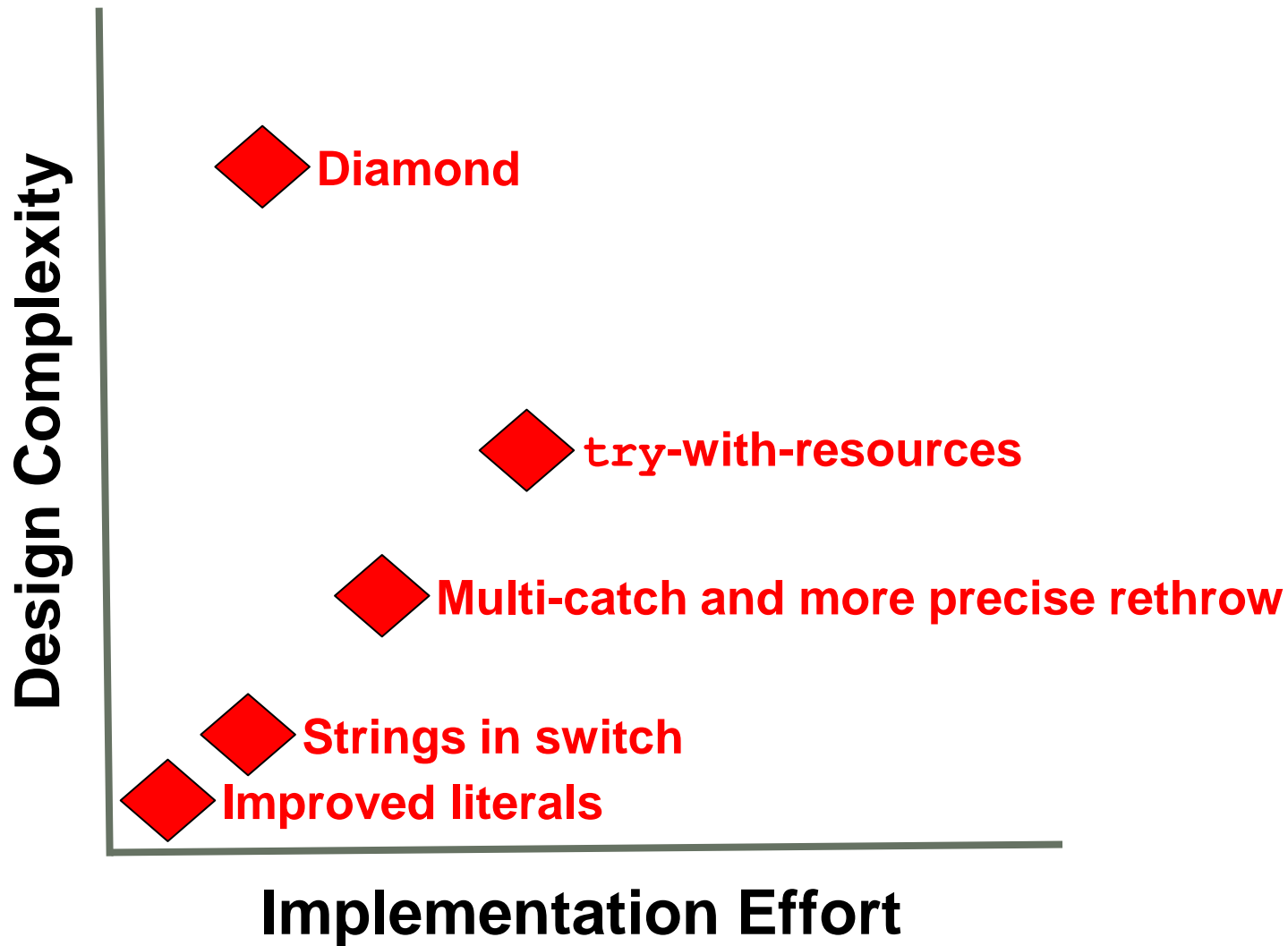
```
class a$1 extends Box<List<capture of ?>> {...}
```

Therefore, disallow non-denotable types in diamond inference.

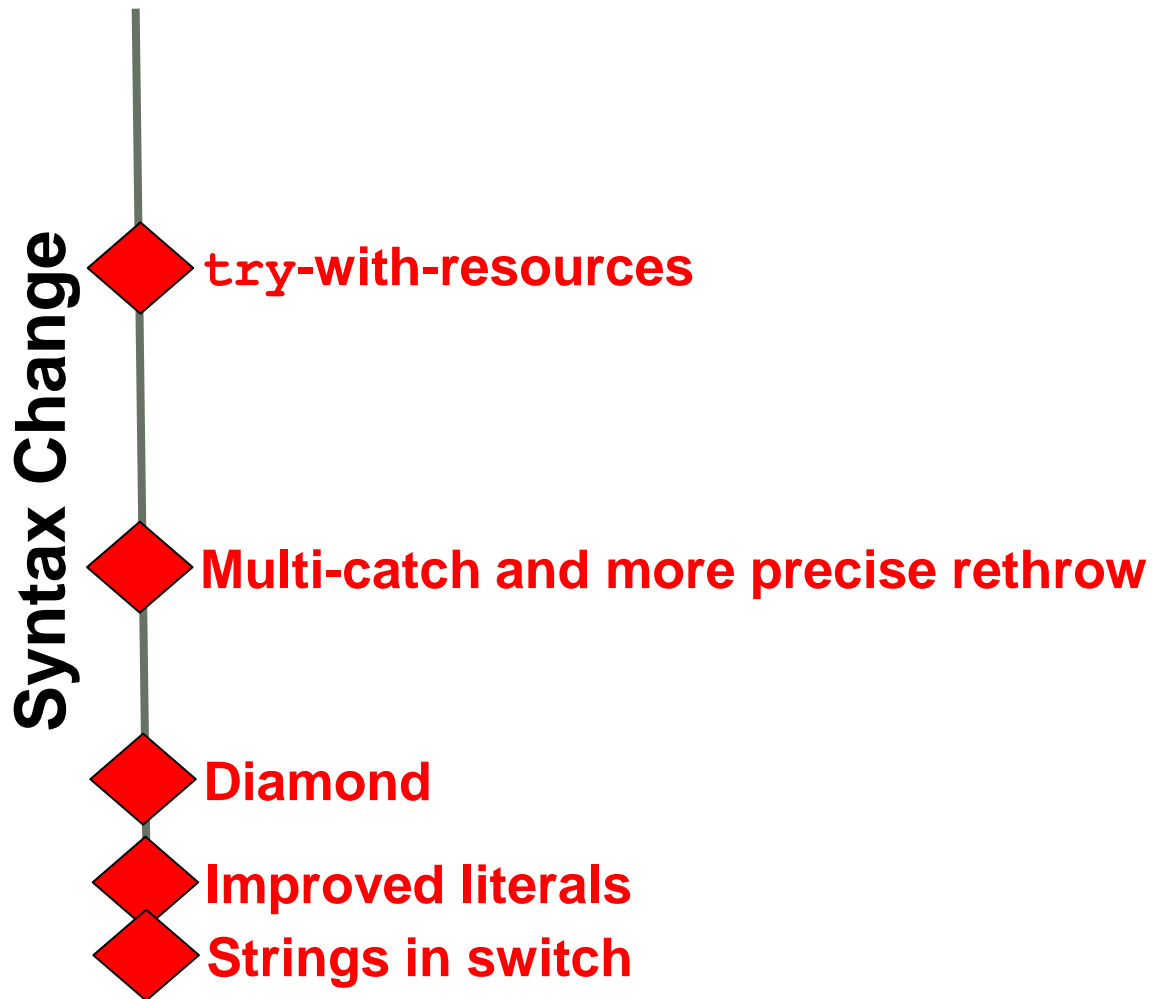
Lesson: keeping the future safe from the past

- Language features over time
 - Anonymous class in JDK 1.1 (1997)
 - Generics in JDK 5 (2004)
 - Diamond in JDK 7 builds (2009)
- Seemingly unrelated features can have deep semantic interactions!

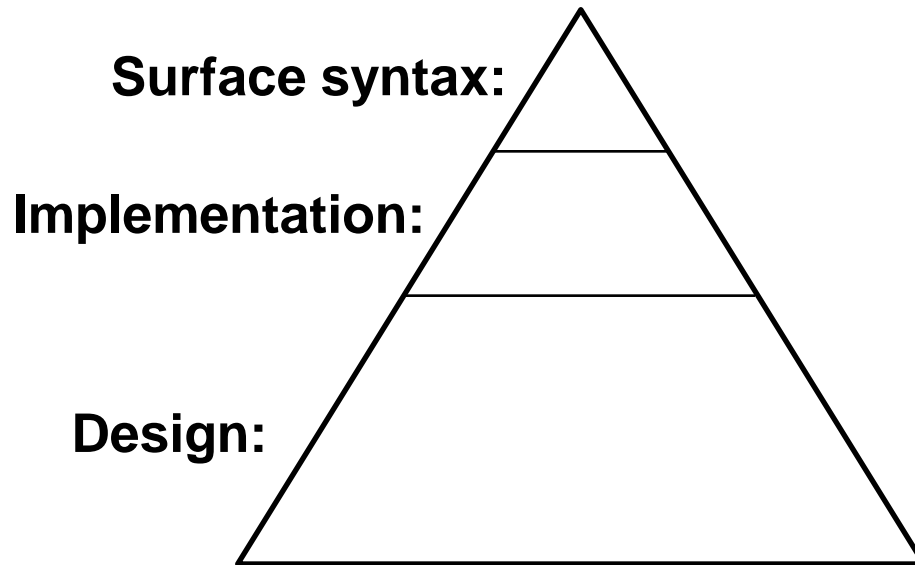
Sizing up the new features



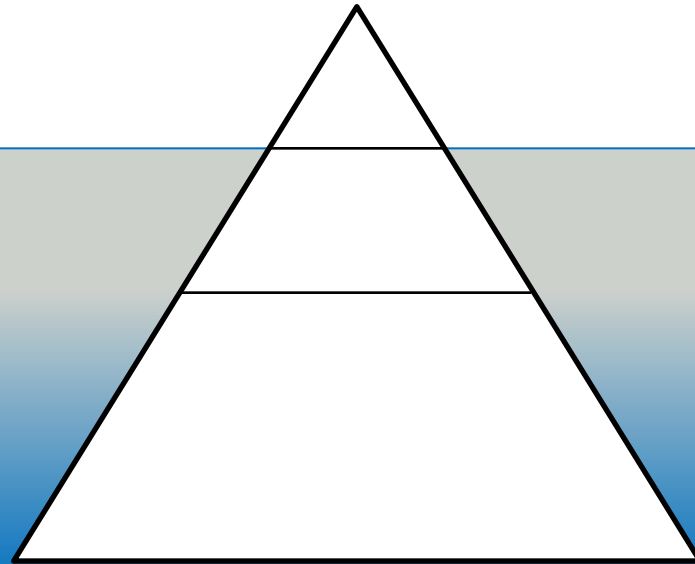
Sizing up the *syntax* of new features



Where does the effort go?



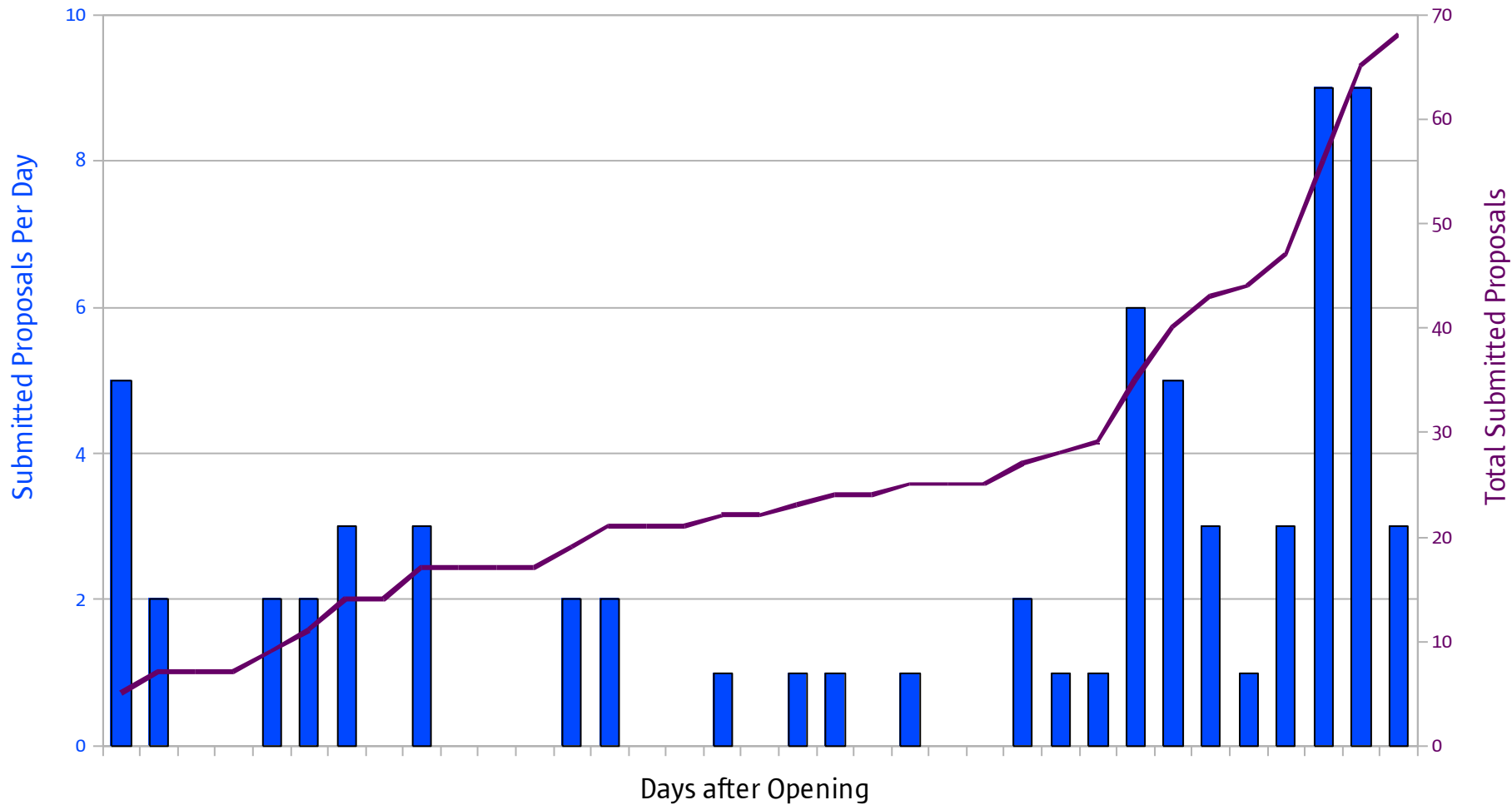
An iceberg!





All in the open

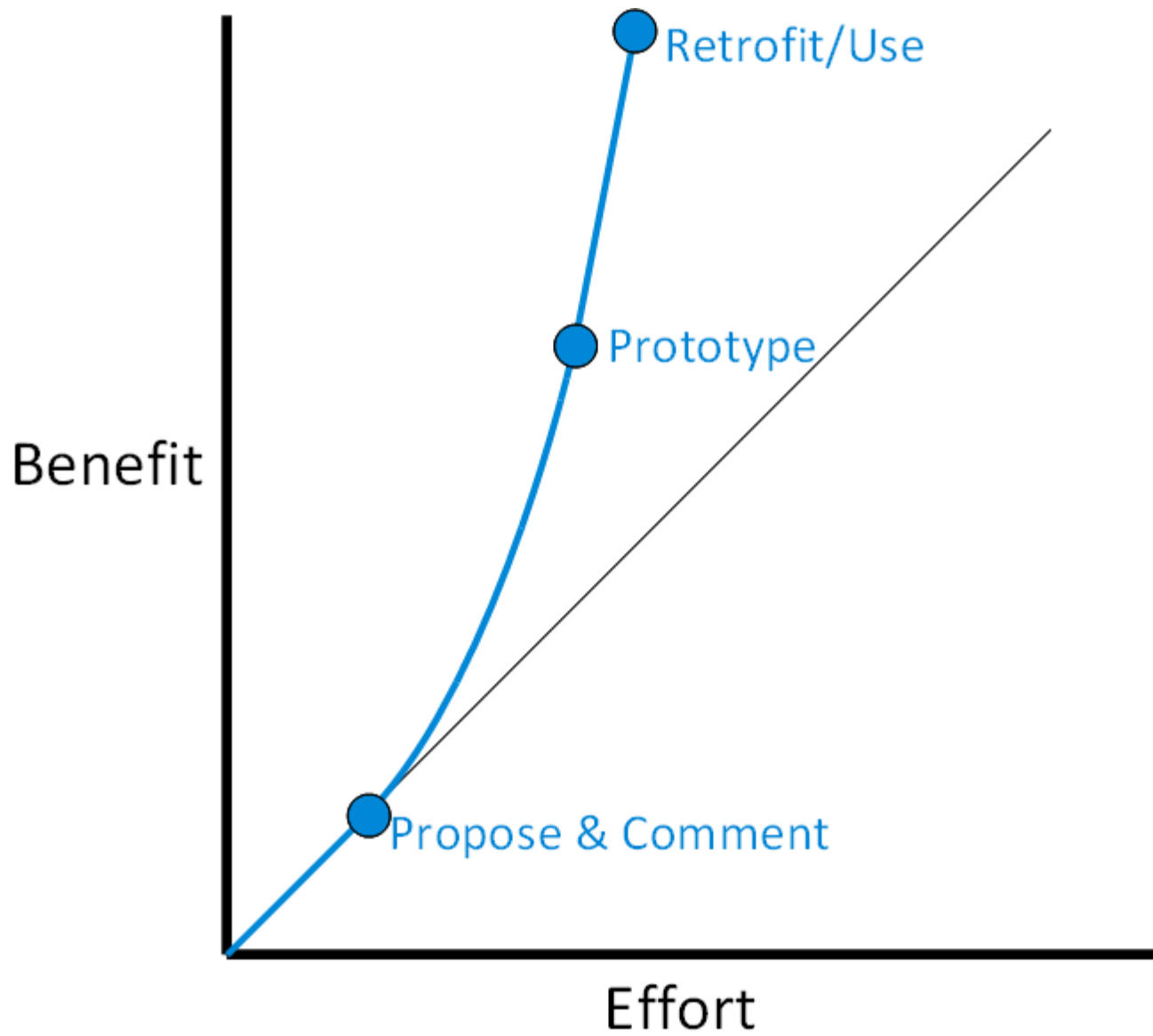
Project Coin Proposal Submissions



coin-dev traffic		Cumulative
February 2009	36	36
March 2009	1288	1324
April 2009	307	1631
May 2009	313	1944
June 2009	130	2074
July 2009	59	2133
August 2009	44	2177
September 2009	81	2258
October 2009	124	2382
November 2009	11	2393

Cumulative non-commit traffic May 2007 to November 2009

hotspot-dev	740
corelibs-dev	626



For the future...

- Extended invitation to participate in *doing the work*
- Lots of time communicating
- Want better signal to noise ratio
 - “Pokka dots on the bikeshed look good!”
- Burden of proof on the proper party
- Will consider a staged process, a proposal can't advance to consideration without a prototype

Summary

- Coin features affect the *bodies* of methods, not their signatures
- Rounding off sharp corners of generics
 - Diamond
 - Varargs warnings
- Increase % of code for non-exceptional circumstances
 - Multi-catch
 - `try-with-resources`
- Consistency and clarity
 - Strings in switch
 - Literal improvements

Conclusions

- Features easier to use than to develop!
- Expect increasing use of quantitative design
- Tooling support important along the way
- Project Coin features
 - Remove superfluous text making programs more *readable*
 - Encourage writing programs that are more *reliable*
 - Play well with past and future changes
- Features ready to try, please give us feedback!



<http://www.jcp.org/en/jsr/summary?id=334>

<http://openjdk.java.net/projects/coin>

<http://jdk7.dev.java.net/>

Q & A

<http://blogs.sun.com/darcy>